

Induction and Recursion

Principle of mathematical induction (weak form):

$$\frac{P(n_0), n_0 \in \mathbb{Z} \\ P(n) \rightarrow P(n+1) \quad \forall n \in \mathbb{Z}, n \geq n_0}{P(n) \quad \forall n \in \mathbb{Z}, n \geq n_0}$$

$P(n)$ in the second line is called the induction hypothesis. Can prove this using the Well-Ordering Principle.

Principle of mathematical induction (strong form):

$$\frac{P(n_0), n_0 \in \mathbb{Z} \\ \forall n \in \mathbb{Z}, n \geq n_0 \quad (P(k) \quad \forall k \in \{ n_0 \dots n \}) \rightarrow P(n+1)}{P(n) \quad \forall n \in \mathbb{Z}, n \geq n_0}$$

Follows from the weak form.

Example: prove that $\forall n \in \mathbb{N}, 8 \mid (3^{2n}-1)$

Definition of the factorial:

$$0! = 1, \quad n! = n(n-1)(n-2)\dots 1.$$

Stirling's approximation: $n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$

Definition: A set A of integers is called an *ideal* iff

$$0 \in A,$$

$$a \in A \rightarrow (-a) \in A,$$

$$a, b \in A \rightarrow a + b \in A$$

prove that $n\mathbb{Z}$ is an ideal

Recursive definitions

Example:

$0!=1, \forall n \in \mathbb{N}, n! = (n-1)! n$ - check Stirling's approximation!

Prove that $n! \geq 3^{n-2}$

Definition:

In this chapter, we define a **sequence** as a function $f: \mathbb{N} \rightarrow \mathbb{R}$.

A sequence can be defined recursively
by **initial condition** and **recursive relation**

example: $a_1=2, a_2=3, a_{k+1} = 2a_k + a_{k-1} \forall k > 2$

arithmetic sequence: $a_1 = a, a_{k+1} = a_k + d$

d is **common difference**

$$S_n = \frac{n}{2} [2a + (n-1)d]$$

geometric sequence: $a_1 = a, a_{k+1} = ra_k$

r is **common difference**

$$S_n = \frac{a(1-r^n)}{1-r}$$

Fibonacci sequence: $f_1 = 1, f_2 = 1, f_{k+1} = f_k + f_{k-1}$

Pseudorandom numbers: $x_{k+1} = (ax_k + b) \pmod{n}$

Next: Solving recurrence relations of the form

$$a_n = ra_{n-1} + sa_{n-2} + f(n)$$

using characteristic polynomials

Theorem: given a *homogeneous* recurrence relation

$$a_n - ra_{n-1} - sa_{n-2} = 0,$$

the solution is

$$a_n = \begin{cases} c_1 x_1^n + c_2 x_2^n & \text{if } x_1 \neq x_2, \\ c_1 x^n + c_2 n x^n & x_1 = x_2 = x, \end{cases} \quad (*)$$

x_1, x_2 are roots of $x^2 - rx - s = 0$

In this case $x^2 - rx - s$ is called *characteristic polynomial*.

The term comes from linear algebra:

$$V_1 = \begin{pmatrix} a_2 \\ a_1 \end{pmatrix}, \quad V_n = \begin{pmatrix} a_n \\ a_{n-1} \end{pmatrix}, \quad A = \begin{pmatrix} r & s \\ 1 & 0 \end{pmatrix}, \quad V_n = AV_{n-1}$$

Looking for a particular homogeneous solution: $V_n = xV_{n-1}$

$$V_n = x^n \xi, \quad A\xi = x\xi \rightarrow \det(A - x) = 0 \rightarrow x^2 - rx - s = 0,$$

$$\xi = \begin{pmatrix} x \\ 1 \end{pmatrix}. \text{ General homogeneous solution is the linear combination}$$

Theorem: The general solution to the recurrence relation

$$a_n - ra_{n-1} - sa_{n-2} = f(n)$$

is the sum of a particular solution and the general solution (*) to the homogeneous recurrence relation.

Definition:

Generating function of a sequence a_0, a_1, a_2, \dots is $a_0 + a_1x + a_2x^2 + \dots$

Examples

$$\frac{1}{1-x} = 1 + x + x^2 + x^3 + x^4 + x^5 + \dots$$

$$\frac{1}{(1-x)^2} = 1 + 2x + 3x^2 + 4x^3 + 5x^4 + 6x^5 + \dots$$

Example: construct generating function for Fibonacci sequence

$$a_0 = 1, a_1 = 1, a_{k+1} = a_k + a_{k-1}$$

$$f(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + \dots$$

$$xf(x) = a_0x + a_1x^2 + a_2x^3 + \dots$$

$$x^2f(x) = a_0x^2 + a_1x^3 + \dots$$

$$\begin{aligned} (1-x-x^2)f(x) &= a_0 + (a_1 - a_0)x + \dots + (a_{k+1} - a_k - a_{k-1})x^{k+1} + \dots \\ &= 1 \end{aligned}$$

$$\rightarrow f(x) = \frac{1}{1-x-x^2}$$

Now we can in principle calculate the coefficients.

Towers of Hanoi

Problem: construct a recursive definition of the algorithm that solves the puzzle for any number of disks

and compute the number of moves

